# HDL IMPLEMENTATION OF DATA COMPRESSION: LZW ALGORITHM

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#### Abstract

This paper presents LZW data compression algorithm which is implemented using finite state machine technique. The proposed algorithm enhances the performance by using less number of bits than their ASCII code, utilizing content addressable memory arrays. Thus the text data can be effectively compressed using the proposed algorithm. Simulation results using Xilinx tool shows an improvement in lossless data compression scheme. In addition to this, the proposed technique results in reduced storage space by 60.25% and increased compression rate by30.3%.

Keywords - Compression rate, English and Binary text.

#### Introduction

Data compression is often referred to as coding, where coding is general term showing any special representation of data which satisfies a given need. Information theory is defined as the study of efficient coding. Data compression may be viewed as a branch of information theory in which the primary objective is to minimize the amount of data to be transmitted. Data compression has an important role in the area of data transmission and data storage. It plays a key role in information technology. The reduction of redundancies in data representation in order to decrease data storage requirement is defined as data compression. It used to less usage of resources such as data space or transmission capacity. Data compression is classified as lossless and lossy compression. Lossless compression is used for text and lossy compression for image.

The first data compression technique "Morse code" was invented in 1838. Morse code was used in telegraph. In 1977, Abraham Lempel and Jacob Ziv suggested the basic idea of pointer-based encoding. In 1980, Terry Welch invented LZW algorithm which became the popular technique for general-purpose compression systems. It was used in programs such as PKZIP as well as in hardware devices. Lempel-Ziv-Welch proposed a variant of LZ78 algorithms, in which compressor never outputs a character, it always outputs a code. To do this, a major change in LZW is to preload the dictionary with all possible symbols that can occur. LZW compression replaces strings of characters with codes.

#### Merits of Data Compression

- 1. It reduces the data storage requirements.
- 2. It also provides rich-quality signals for audio data representation.
- 3. Data security can also be greatly enhanced by encrypting the decoding parameters and transmitting them separately from the compressed database files.
- 4. Data compression obviously reduces the cost of backup and recovery of data in computers system by storing large back-up database files in compressed form.

#### Data Compression Model

The block diagram of data compression model is described in figure 1.

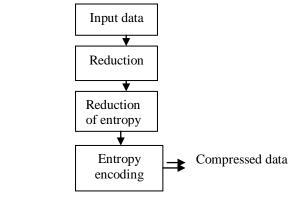


Figure 1: Data compression model

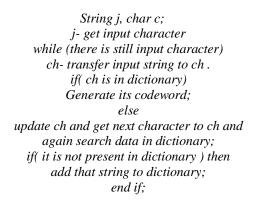
A data compression model consists of three major stages which are redundancy, reduction in entropy and entropy encoding.

# Data Compression Algorithm: LZW (Lempel- Ziv Welch) Algorithm

There are many algorithms which have been used for data compression like Huffman and Lempel-Ziv-Welch (LZW), arithmetic coding. LZW algorithm is the most popular algorithm. LZW algorithm is just like a greedy approach and divides text into substrings. Like the LZW algorithm proposed in [2] LZW algorithm has both compression and decompression techniques which is explained as

#### LZW Compression Algorithm

LZW compression algorithm is dictionary based algorithm which always output a code for a character. Each character has a code and index number in dictionary. Input data which we want to compress is read from file. Initially data is entered in buffer for searching in dictionary to generate its code. If there is no matching character found in dictionary. Then it will be entered as new character in dictionary and assign a code. If Character is in dictionary then its code will be generate. Output codes have less number of bits than input data. This technique is useful for both graphics images and digitized voice



Compression example: consider a string "BAABAABB" is given to LZW algorithm. Figure 2 shows the steps done by LZW to generate the output code is "1211211C". In following example when input string (BAABAABBC) is given as a text to LZW compression algorithm. Initially every single character will save in buffer. When 'B'is move to buffer "parse string" then it will replace by 1.Character has its own ASCII code of 7 bit. In case of B, it has 65 as ASCII code. But in dictionary it will replace by 1.So, less number of bits will be used to represent character. Similarly, AA will move forward and generating its code which is also fewer bits than original. BAA is saved in buffer its code is generated from both AA and B's code words that is defined as 12. At last

when full string has been searched in dictionary then its output will be generated as 1211211C.

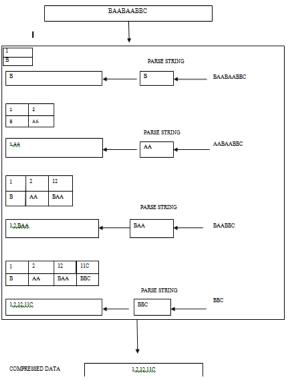


Figure 2: Example of LZW algorithm

#### LZW Decompression Algorithm

In LZW decompression algorithm, it needs to take the stream of code output from the compression algorithm, and use them to exactly recreate the input stream. Decompression algorithm is shown as:

ch = output code while (there is still data to read) code =get input character; if (code is not in the dictionary) entry =get translation of code; else entry=get translation of output code; output entry; ch =first character in entry add output code + c to the dictionary output code = code; In decompression algorithm, code will be searched in dictionary and its character will be output.

# Implementation of LZW Algorithm

#### Implementation of LZW Algorithm

The proposed finite state machine diagram of LZW algorithm is shown in figure 3.

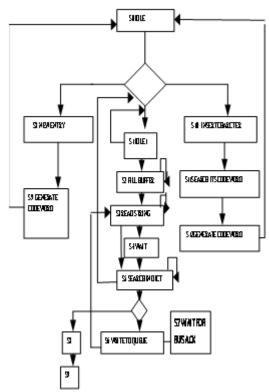


Figure 3: Finite state machine Diagram of LZW algorithm

LZW algorithm initially has idle state. New character has been added to dictionary when no longer match will found in search process. LZW algorithm is execute state S8 for performing adding operation in dictionary. Dictionary is based on content access memory technique which has both content as well as code in it.

Decompressor has reviewed same process since it is possible to have input codes for searching in dictionary to recreate its original string. Individual character's code can be also viewed in dictionary

Table 1: Specifications	of FSM state for LZW Algorithm
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State	Description	
S0 idle	Initial state reset the system	
S1 idle1	Initialization of signal	
S2Fill buffer	Transfer text from file to buffer	
S3Read string	Read character by character for searching	
S4 wait	For waiting	
S5 search in dict	For searching in dictionary by signal character	
S6Write to queue	Save output to output buffer	
S7 wait for ack	Wait for Bus acknowledge	
S8New Entry	Adding new entry	
S9Generate codeword	To generate codes	
S10	Insert single character	
S11search its codeword	Check in dictionary	
S12 generate codeword	Display codeword	
Decompression	For performing decompression	

# Improvement of the Dictionary Storage Method

LZW algorithm is mainly used for compressing character but not numeric. Every character has ASCII code which is of 7 bits. But in our proposed algorithm we have to replace character with 5 bit code in dictionary to improving data compression rate.

# **Experimental Results**

LZW Compression algorithm is modelled in VHDL. The syntax of the RTL design is checked by using Xilinx tool.

#### Simulation Results

In the proposed work, the simulations results are done using Xilinx ISE Simulator. Simulation results show an improvement in lossless data compression scheme. In addition to this, the proposed technique results in reduced storage space by 60.25% and increased compression rate by 30.3%.

#### LZW Compressor Result

Figure 4 shows that input is given to LZW Compressor through text file. "Connect the input to logic one & two & three++\*"string is entered to it. Input string having collection of special characters, alphabets. Whole text will transfer to buffer "data read" when data\_write=1, load=1, clear=1.Rd\_b=0, wr\_b=1, data\_write=0 and lzw\_search=1 are given to start searching process to find longest match in content access memory arrays. There are two main counters which are used for searching process. "Count1" is used for searching character in dictionary. If character is present in dictionary then its code is saved in other buffer that is "del1". "Count" buffer is shifted to next value and start to point next character present in input data. All searched characters will save in "check" buffer. Once the content of check buffer is equal to content of "data read" buffer then searching process indicate to completed and their codes will save in "del1" buffer which is shown in figure5 compressed output is generated through file shown in figure 7.

Input text – connect the input to logic one & two & three. Output text-12114514615730171425142515

] 1- Notepad	
File Edit Format View Help	
connect the inputs to logic one & two & three++&*.	Å

Figure 4: Enter data through file which we want to compress

Given Input text – connect the input to logic one & two & three.

Name	Value		211.ps	212.ps	213 ps	214ps	215 ps	215 ps	217ps	218 ps	219 ps 22
le bus_ack	U										
la match found	U U										
le buffer0_bus_data											
<ul> <li>M q7:0</li> </ul>	-						uuuu				
• • • • • • • • • • • • • • • • • • •	0000000										
de11:50	[0,10101,10,11,11,			00301.30.11.11.1.3							
tag[0:255]	[0000000,0000000										
data[0:255]	(0000000,0000000	[JUUUUUU	,			,			,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,,		ստուստութ.
count_1	0						0				
kkey[1:50]	['c','o','n','n','	1		[c,o,n,n,e,c,t,	,t,h,e, ,i,n,p,	u,t,s, ,t,o, ,1,		,a, ,t,w,o, ,a	, ,t,h,r,e,e,+,	+,&, ,.]	
e kkey_ink	0						0				
🕨 幡 check[1:50]	['c','o','n','n',	[c,o,n,n	,e,c,t, ,t,h,e,	,ï,n,p,u,ť,s, 🌶	առառառա	LNUNUNUNU	λαναναγικ	μλαλαγά	(MINININI	NUNUNUN	ແກແກແກແກ
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le dock	0										
le endoffile	1										
endoffile1	0										
dataread[1:50]	['c','o','n','n','			[c,o,n,n,e,c,t,	,t,h,e, ,ï,n,p,	u,t,s, ,t,o, ,T,	,g,T,C, ,o,N,e,	,a,,t,w,o,,a	, ,t,h,r,e,e,+,	+, 8, •, .]	
lie linenumber	10011						10011				
Le linenumber1	0						0				
16 .	10100						10100				
le count decom	1						1				
Le count	10100						10100				
Up counti	100	10	11	100	101	110	111	1000	1001	1010	1011
Se country											

Figure 5: Searching process (Searching each character from dictionary)

Simulation for LZW Compression algorithm observed on Xilinx tool. When 350 bits entered to LZW compression algorithm then it is transmitted to 119 bits and clock rate for simulation is 493 ps.

	bus_ack	U										
18	match_found	U										
1.2	buffer0_bus_data	1										
- D- 🎫	q[7:0]	00000000						30000000				
	q_tmp[7:0]	00000000						JUUUUUUU				
p- 🞫	de1[1:50]	[0,10101,10,11,11,	[0, 10	101, 10, 11, 11, 1, 1010	1,0,101,1110,110,0	101,111,11,0,1000	,1110,0,101,1110,0	,101,1111,10,10111	, 111,0, 101, 10, 11,0,	0,0,101,1110,1000,	0,0,0,101,1110,110,	100, 1, 0, 0, 0, 0, 0]
	tag[0:255]	(00000000,00000000	[ບບບບບບບບ,	,00000000,00000	100,00000000,000	00000,00000000,				00000000,000000		00000,0000000,0
p- ==/	data[0:255]	100000000,000000000	ບບບບບບບບ	0000000,00000	00,0000000,000	00000,00000000,	0000000,00000	00,0000000,000	00000,00000000,	00000000,00000		00000,000000000
18	count_1	0						0				
p- =	kkey[1:50]	['e','e','n','n','			['c','o','n','n','e','c','t','	','t','h','e',' ','l','n','p','	u','t','s','','t','o',' ',T,'e	','g','i','c',' ','o','n','e',	','&', ','t','w','o', ','&	,'','t','h','r','e','e','+','	+','&','*','.']	
18	kkey_ink	0						0				
p- 🞫	check[1:50]	['e','o','n','n','			c','o','n','n','e','c','t',' '	'ť,'h','e',' ','i','n','p','u	,'t','s', ','t','o', ',T,'o'	'g','','c', '','o','n','e','	'&', ','t','w','o', ','&','	','t','h','r','e','e','+','+	','&','*',NUL]	
1.2	reset1	0										
18		NUL						NUL				
	clock	0										
- 18	endoffile	1										
18	endoffile1	0										
- IF 🤜	dataread[1:50]	['e','o','n','n','			['c','o','n','n','e','c','t','	','t','h','e',' ','i','n','p','	u','t','s','','t','o','','T,'e	','g','i','c',' ','o','n','e',	','&', ','t','w','o', ',&	,'','t','h','r','e','e','+','	+','&','*','.']	
	linenumber	110001						110001				
18	linenumber1	0						0				
	0	110010						110010				
18	count_decom	1						1				
18	count	110010						110010				
18	count1	110	100	101	110	(111)	1000	1001	1010	1011	1100	1101

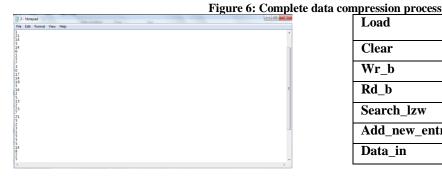


Figure 7: Compressed output generate on file

Output text-12114514615730171425142515

#### **RTL** view of LZW Compressor

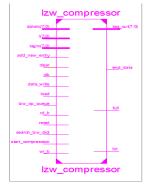


Figure-8: RTL view of LZW Compressor

This RTL view shows the signals which are used for proposed LZW data compression algorithm. Reset, clock, start compression used for initialization of data compression. Load, data write, wr b, rd b are signals used for buffer in LZW algorithm.

Search\_lzw is for searching data in dictionary. The signal description of this proposed algorithm is shown in table 2.

#### Signal description of LZW Compressor

Table 3: Input/output signals with Remarks
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Name	Description
Reset	To reset
Clock	Provide clock
Start_compression	Signal for start compres- sion
Data_write	Signal for write data

inpression process	
Load	For data load in buffer
Clear	Clear buffer
Wr_b	Signal for write and read
Rd_b	Signal for write and read
Search_lzw	For searching
Add_new_entry	For adding new data
Data_in	Enter value

#### Verification and Synthesis

For system verification, we successfully execute proposed LZW algorithm. Test case for finite state machine is generated in VHDL. The synthesis result of LZW compression algorithm is summarized in table 4. The synthesis report shows device utilization summary.

Table 4: Device U	Itilization Summary
Number of Slices	3606 out of 6144
	58%
Number of Slice Flip	4097 out of 12288
Flops	33%
Number of 4 input LUTs	4190 out of 1288
	34%
Number of IOs:	30
Number of bonded IOBs:	30 out of 240
	12%
IOB Flip Flops:	1

# Table 4. Design Littlenstion C

# Conclusion

In order to get better compression rate. The proposed dictionary based LZW algorithm can replace their codes with 5 bits instead of 7 bits ASCII code. The proposed LZW algorithm is evaluated by finite state machine technique. With this technique we have observed that storage space is reduced up to 60.25% and compression rate improved up to 30.3%. We analyze compression rate with different number of input bits on Xilinx tool.

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# Biographies

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